97-0617: Worst Case Buffer Requirements for TCP over ABR

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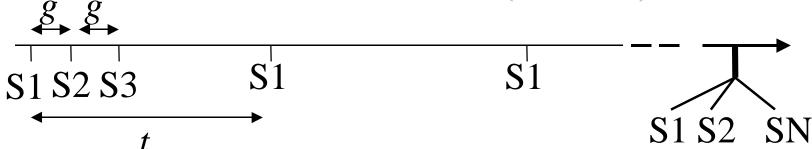


- Generating Worst Case TCP Traffic
- Analytical Results
- □ Simulation Results
 - □ Effect of number of sources
 - □ Effect of RTT, MSS, and other parameters

Worst Case TCP Traffic

- □ Sources can retain high ACR, if they send packets within 500 ms.
- Many such sources with high ACR can dump a large amount of data
- □ Worst case is when all the sources dump the maximum window size

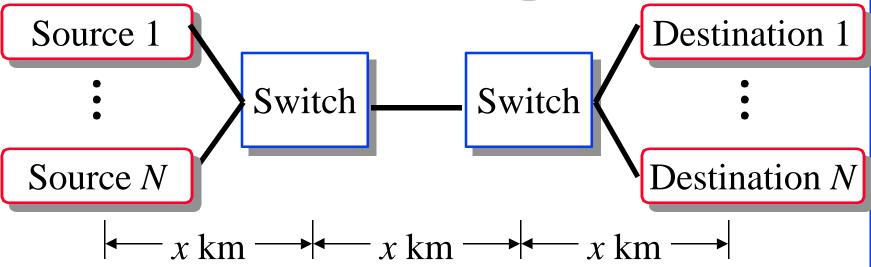
Worst Case (Cont)



- Each source sends one packet every 't' milliseconds. t < 500 ms.
- After several packets, the congestion window reaches the maximum for each source
- □ Sources synchronize and dump large burst at the same time.
- □ To avoid overload initially, the sources are uniformly spaced \triangleright *k*th source sends its first packet at ' $k \times g$ ' μ s.

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N-Source Configuration



- □ All links 149.76 Mbps. Lengths x = 2000, 1000 km
- All traffic unidirectional. Worst case TCP traffic
- □ Parameters: # of sources={2, 3, 5, 10, 20, 30, ..., 200} Infinite buffer size.

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Analytical results

- Buffer requirement is reflected in maximum switch queue size.
- □ Let cwnd_max = Max congestion window of TCP
- \square When $N \leq \lfloor t/g \rfloor$
 - □ Max Q length \approx N × cwnd_max/48 (formula 1)
- □ When $N > \lfloor t/g \rfloor$

48 bytes/cell (formula 2)

- □ Max Q length = N*PCR*t (PCR is peak cell rate)
- Queue length is given in terms of number of cells

Analytical results (Cont)

- With few sources, switch does not get congested even when sources reach their maximum window, ACRs can be high. Formula 1 applies here.
- With many sources, switch detects congestion and gives feedback. ACRs are low. Formula 2 applies here.

Simulation Parameters

□ Source: Parameters selected to maximize ACR

$$TBE = 512$$

$$CDF(XDF) = 0.5$$

$$ICR = 10 \text{ Mbps}$$

$$CRM (Xrm) = [TBE/Nrm]$$

$$ADTF = 0.5 \text{ sec}$$

$$PCR = 149.76 \text{ Mbps}, MCR = 0, RIF (AIR) = 1,$$

$$Nrm = 32$$
, $Mrm = 2$, $RDF = 1/512$, $Trm = 100ms$,

$$TCR = 10 \text{ c/s}$$

- □ Traffic: TCP/IP with worst case traffic
- □ Switch: ERICA+

Averaging interval = $min\{100 \text{ cells}, 1000 \mu s\}$

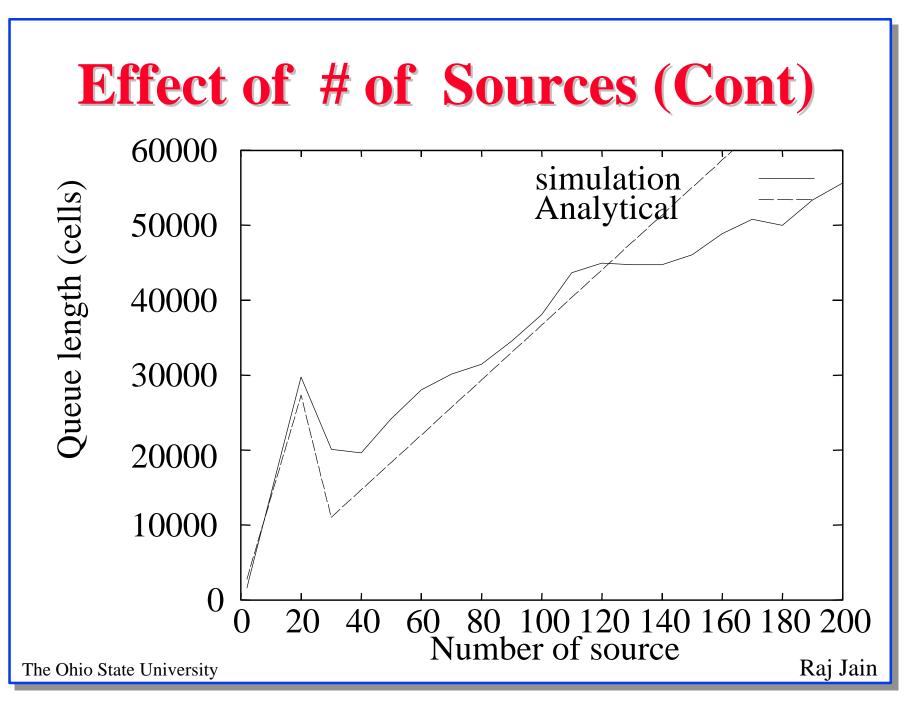
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Effect of Number of Sources

# TCP	Q Size (Cells)				
Srcs	Simul.	Analyt.			
2	1575	2730			
3	3149	4095			
5	6297	6825			
10	14131	13650			
20	29751	27300			
30	20068	11010			
40	19619	14680			
50	24162	18350			
60	28006	22020			

# TCP	Q Size (Cells)				
Srcs	Simul.	Analyt.			
70	30109	25690			
80	31439	29360			
90	34530	33030			
100	38088	36700			
120	44939	44040			
140	44744	51380			
160	48880	58720			
180	49961	66060			
200	55618	73400			

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Effect of # of Rources (Cont)

- □ Analytical results: For t = 1 ms, g = 50 µs, MSS = 512 bytes, cwnd_max = 64 kB
 - Q = N*1365

for $N \le 20$ (formula 1)

Q = N*367

for N > 20 (formula 2)

- ☐ The zig-zag shape is due to the two formulas
- □ The simulation agrees well with the analytical results for $N \le 20$.
- □ The maximum queues occurred at predicted times (details in the contribution)

Effect of # of Sources (Cont)

- Buffer size increases linearly as number of sources increase
- □ As N increases, load increases
 - \Rightarrow ERICA+ controls the queue lengths \Rightarrow Less than analytical queue lengths

Sensitivity Analysis

#	mss/g/t/d	N=3	N=10	N=30	N=40	N=50	N=100
1	512/50/1/1000	3171	14273	20068	19619	24162	35687
2	512/50/1/2000	3171	14273	19906	27567	30872	75083
3	512/50/10/1000	3172	14274	45994	61854	77714	150453
4	512/50/10/2000	3172	14274	45994	61854	77714	150458
5	512/100/1/1000	3171	14273	19283	20080	24164	NA
6	512/100/1/2000	3171	14273	21241	32314	35961	NA
7	512/100/10/1000	3172	14274	45994	61854	77714	NA
8	512/100/10/2000	3172	14274	45994	61854	77714	NA
9	1024/50/1/1000	3040	13680	18650	18824	23542	NA
10	1024/50/1/2000	1542	5612	19131	22934	29163	NA
11	1024/50/10/1000	3040	13680	44080	59280	74480	NA
12	1024/50/10/2000	3041	13681	44081	59281	74481	NA
13	1024/100/1/1000	3040	13680	18591	19600	24314	NA
14	1024/100/1/2000	1403	5556	17471	24412	30533	NA
15	1024/100/10/1000	3040	13680	44080	59280	74480	NA
16	1024/100/10/2000	3041	13681	44081	59281	74481	NA Rai Jair

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Sensitivity Analysis: Results

- MSS = 512, 1024 bytes, t = 1, 10 ms, g = 50, 100 µs, Link distance = 1000, 2000 km Two values for each of the 4 parameters \Rightarrow 16 experiments.
- Segment size does not affect queue sizes
- □ If the network is not overloaded then round trip time has no effect (Expt. 3 and 4)
- □ If the network is overloaded then a larger round trip gives larger queue lengths (Expt. 1, 2 for N = 30, 40, 50)



- ☐ In worst case, the buffer requirements depend on the number of sources, network congestion status (overloaded or underloaded) and round trip time
- □ It is not affected by maximum segment size.