Transport Protocols

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These slides are available on-line at:

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q TCP

- q Key features
- q Header format
- q Mechanisms
- q Implementation choices
- q Slow start congestion avoidance



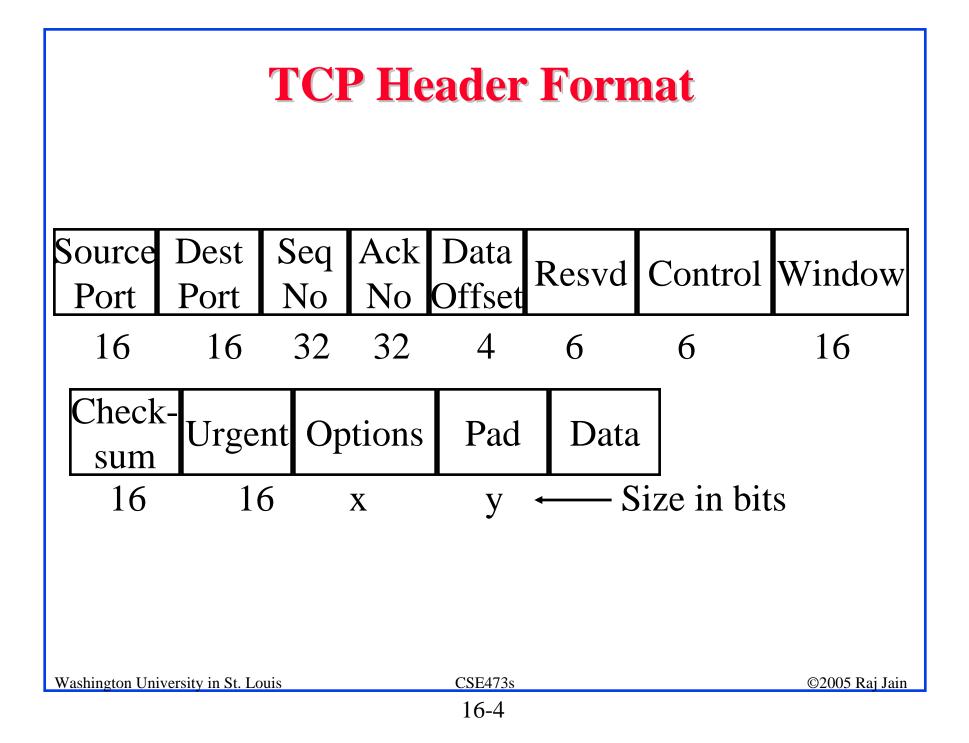
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TCP

- q Transmission Control Protocol
- **q** Key Services:
 - q Send: Please send when convenient
 - q Data stream push: Please send it all now, if possible.
 - q Urgent data signaling: Destination TCP! please give this urgent data to the user
 (Urgent data is delivered in sequence. Push at the source should be explicit if needed.)
 - q Note: Push has no effect on delivery.Urgent requests quick delivery

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TCP Header

- q Source Port (16 bits): Identifies source user process
- q Destination Port (16 bits) 20 = FTP, 23 = Telnet, 53 = DNS, 80 = HTTP, ...
- q Sequence Number (32 bits): Sequence number of the first byte in the segment. If SYN is present, this is the initial sequence number (ISN) and the first data byte is ISN+1.
- q Ack number (32 bits): Next byte expected
- q Data offset (4 bits): Number of 32-bit words in the header
- q Reserved (6 bits)

TCP Header (Cont)

 q Control (6 bits): Urgent pointer field significant, Ack field significant, Push function, Reset the connection, Synchronize the sequence numbers, No more data from sender

URG ACK	PSH	RST	SYN	FIN	
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q Window (16 bits): Will accept [Ack] to [Ack]+[window]-1

TCP Header (Cont)

- q Checksum (16 bits): covers the segment plus a pseudo header. Includes the following fields from IP header: source and dest adr, protocol, segment length. Protects from IP misdelivery.
- q Urgent pointer (16 bits): Points to the byte following urgent data. Lets receiver know how much data it should deliver right away.
- **q** Options (variable):

Max segment size (does not include TCP header, default 536 bytes), Window scale factor, Selective Ack permitted, Timestamp, No-Op, End-of-options

TCP Options

Kind	Length	Meaning
0	1	End of Valid options in header
1	1	No-op
2	4	Maximum Segment Size
3	3	Window Scale Factor
8	10	Timestamp

- q End of Options: Stop looking for further option
- q No-op: Ignore this byte. Used to align the next option on a 4byte word boundary
- q Max Segment Size (MSS): Does <u>not</u> include TCP header

TCP Checksum

- q Checksum is the 16-bit one's complement of the one's complement sum of a pseudo header of information from the IP header, the TCP header, and the data, padded with zero octets at the end (if necessary) to make a multiple of two octets.
- **q** Checksum field is filled with zeros initially
- **q** TCP length (in octet) is not transmitted but used in calculations.

q Efficient in	Efficient implementation in RFC1071.						
Source Adr	Dest. Adr	Zeros	Proto	ocol TCP Len	gth		
32	32	8	8	16			
	TCP Header	r TCP	data				
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16-9							

1's Complement

2's Complement: -ve of a number is 1+complement

- q 1 = 0001 -1 = 1111
- q 2 = 0010 -2 = 1110
- q 3 = 0011 -3 = 1101

1's complement: -ve of a number is it's complement

- q 1 = 0001 -1 = 1110
- q 2 = 0010 -2 = 1101
- q 3 = 0011 -3 = 1100

2's Complement sum: Add with carry

1's complement sum: Add. Add the carry back to the sum

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q 8+9 = 1000 + 1001 = 1\ 0001 => 0001 + 1 = 0010
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Complement of 1's complement sum: 1101

Why: 1's complement sum is independent of the Endian-ness of the machines.

Little Endian = LSB is the left most bit.

Big Endian = MSB is the left most bit

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TCP Service Requests

- q Unspecified passive open:Listen for connection requests from any user (port)
- q Full passive open: Listen for connection requests from specified port
- q Active open: Request connection
- q Active open with data: Request connection and transmit data
- q Send: Send data
- q Allocate: Issue incremental allocation for receive data
- **q** Close: Close the connection gracefully
- q Abort: Close the connection abruptly
- q Status: Report connection status

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TCP Service Responses

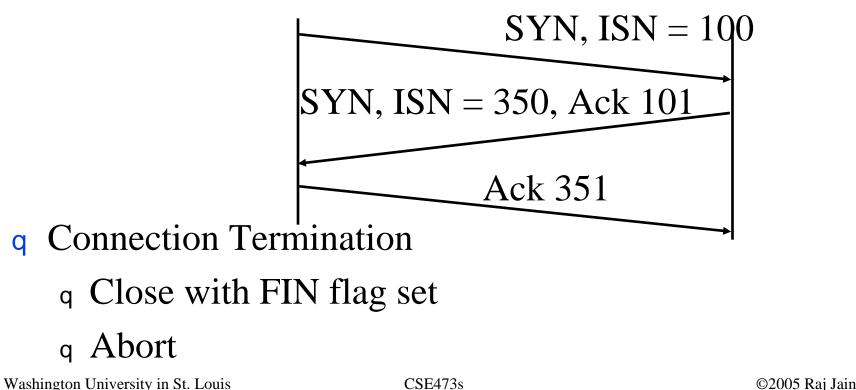
- q Open ID: Informs the name assigned to the pending request
- q Open Failure: Your open request failed
- q Open Success: Your open request succeeded
- q Deliver: Reports arrival of data
- q Closing: Remote TCP has issued a close request
- **q** Terminate: Connection has been terminated
- **q** Status Response: Here is the connection status
- q Error: Reports service request or internal error

TCP Mechanisms

- **Connection Establishment** a
 - q Three way handshake

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q SYN flag set \Rightarrow Request for connection



Data Transfer

- **q** Stream: Every byte is numbered modulo 2^{32} .
- q Header contains the sequence number of the first byte
- **q** Flow control: Credit = number of bytes
- q Data transmitted at intervals determined by TCP Push \Rightarrow Send now
- q Urgent: Send this data in ordinary data stream with urgent pointer
- q If TPDU not intended for this connection is received,the "reset" flag is set in the outgoing segment

Implementation Policies (Choices)

q Send Policy:

Too little \Rightarrow More overhead. Too large \Rightarrow Delay Push \Rightarrow Send now, if possible.

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q Delivery Policy:
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May store or deliver each in-order segment. Urgent \Rightarrow Deliver now, if possible.

q Accept Policy:

May or May not discard out-of-order segments

Implementation Policies (Cont)

- q Retransmit Policy:
 - First only
 - Retransmit all
 - Retransmit individual
 - (maintain separate timer for each segment)
- q Ack Policy:
 - Immediate (no piggybacking) Cumulative (wait for outgoing data or timeout)

Slow Start Flow Control

- **q** Window = Flow Control Avoids receiver overrun
- **q** Need congestion control to avoid network overrun
- q The sender maintains two windows:
 Credits from the receiver
 Congestion window from the network
 Congestion window is always less than the receiver window
- q Starts with a congestion window (CWND) of 1 segment (one max segment size)
 - \Rightarrow Do not disturb existing connections too much.

q Increase CWND by 1 every time an ack is received Washington University in St. Louis CSE473s ©2005 Raj Jain

Slow Start (Cont)

If segments lost, remember slow start threshold q (SSThresh) to CWND/2 Set CWND to 1 Increment by 1 per ack until SSthresh Increment by 1/CWND per ack afterwards **Receiver Window** Congestion Timeout Idle Window SSThresh Interval CWND

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Time

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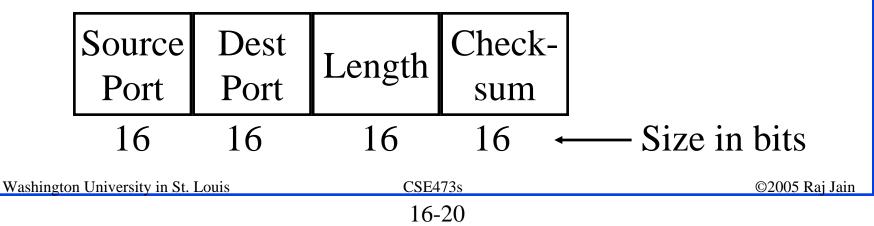
Slow Start (Cont)

- **q** At the beginning, SSThresh = Receiver window
- q After a long idle period (exceeding one round-trip time), reset the congestion window to one.
- q Exponential growth phase is also known as "Slow start" phase
- q The linear growth phase is known as "congestion avoidance phase"

User Datagram Protocol (UDP)

- q Connectionless end-to-end service
- q No flow control. No error recovery (no acks)
- q Provides port addressing
- q Error detection (Checksum) optional. Applies to pseudo-header (same as TCP) and UDP segment. If not used, it is set to zero.

q Used by network management





- q TCP provides reliable full-duplex connections.
- q TCP Streams, credit flow control
- q Slow-start
- q UDP is connectionless and simple.No flow/error control. Has error detection.

Reading Assignment

q Read Chapter 20 of Stallings' 7th edition

Homework

- q Submit answer to Exercise 20.20This homework is worth 30 points.
- **exercise 20.20**: A TCP entity opens a connection and uses slow start. Approximately how many round-trip times are required before TCP can send N segments.
- q Hint: Write down what the CWND and total segments will be after 1 round trips, 2 round trips, 3 round trips,