IP Over SONET

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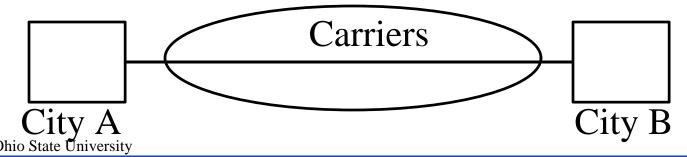
- □ IP over SONET: Trends, Users, Why?
- □ SONET: Key features
- PPP: Key features
- □ SONET vs ATM
- □ IP over SONET: Key Issues
- Products

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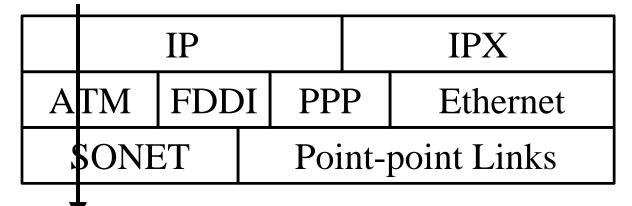
What is SONET?

- Synchronous optical network
- Standard for digital optical transmission (bit pipe)
- Developed originally by Bellcore.
 Standardized by ANSI T1X1
 Standardized by CCITT
 - ⇒ Synchronous Digital Hierarchy (SDH)
- You can lease a SONET connection from carriers

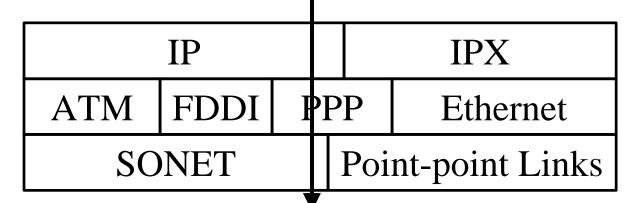


Changing Trends

□ View Until Early 1996:



□ View in Late 1996:



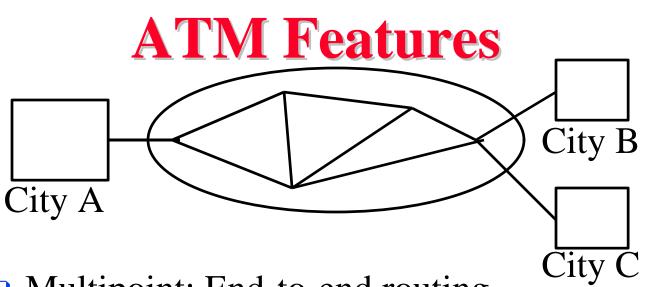
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Trends (Cont)

- Originally, ATM has been designed for high-speed transfer of data, voice, video
- □ Carriers were expected to move to ATM networks
- SONET was designed as a high-speed physical layer for transmission over fiber-optic links
- □ ATM was expected to run over carrier's SONET links
- "IP over SONET" allows IP datagram
 transfers over high-speed carrier links using PPP
- □ SONET is appearing as a competition to ATM

IP over SONET: Users

- □ Trans-Atlantic IP-over-SONET between NYC and Stockholm in Sept'96.
- □ OC-3c SONET ring between US, England, and France by Sprint using CISCO 750's Oct'96.
- □ AT&T and Kokusai Denshin Denwa (KDD)155 Mbps between San Francisco and Tokyo
- □ Internet Servers Inc (ISI) uses for high-volume Web server. Netscape and Yahoo use ISI.
- Ascend and IXC Communications trialed 155 Mbps between Fort Worth TX and El Paso TX (11/96).
- BBN PLANET, UUNET Technologies, and Sprint were experimenting in Jul'96.

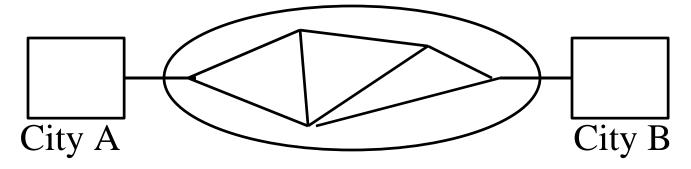


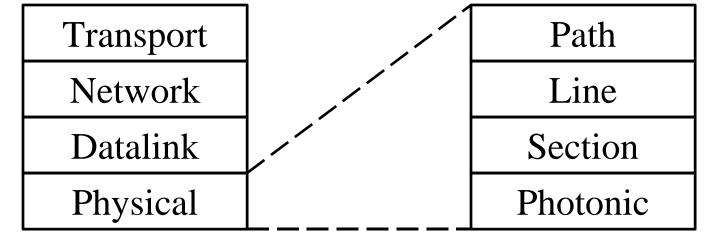
- □ Multipoint: End-to-end routing
- High-speed switching vs slow routing
- Quality of service
- Multiplexing of different QoS: Voice, Video, and data
- □ Signaling: Dynamic bandwidth
- Traffic Management: Overload/underload Control

All features cost overhead.

SONET Protocols

Synchronous Optical Network

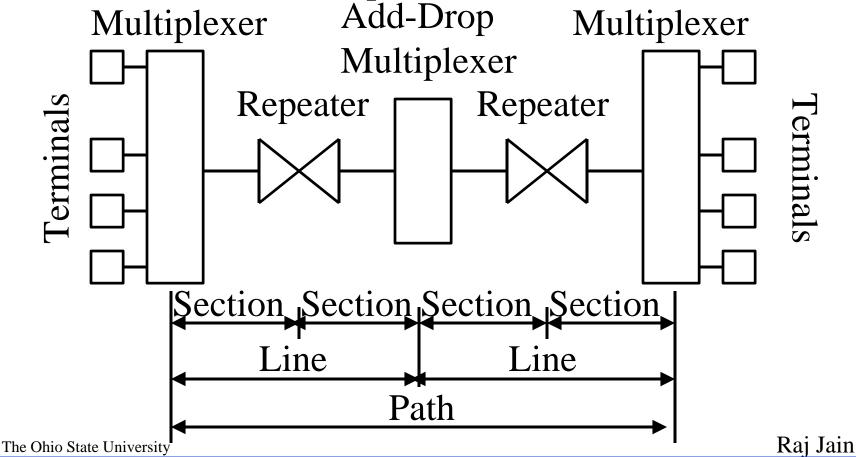




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Physical Components

- → Section = Single run of fiber
- □ Line = Between multiplexers



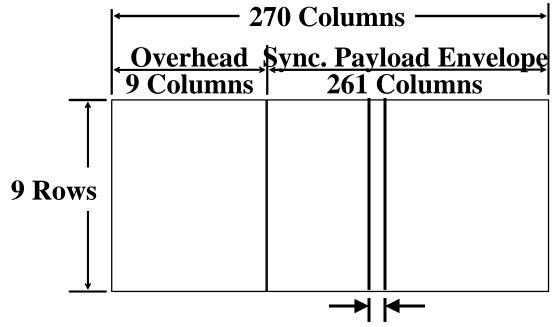
Signal Hierarchy

Synchronous Transport Signal Level $n = STS-n = n \times 51.84$ Mbps STM=Synchronous Transport Module, OC=Optical Carrier level

ANSI	Optical	CCITT	Data Rate	Payload Rate
Designation	Signal	Designation	(Mbps)	(Mbps)
STS-1	OC-1		51.84	50.112
STS-3	OC-3	STM-1	155.52	150.336
STS-9	OC-9	STM-3	466.56	451.008
STS-12	OC-12	STM-4	622.08	601.344
STS-18	OC-18	STM-6	933.12	902.016
STS-24	OC-24	STM-8	1244.16	1202.688
STS-36	OC-36	STM-12	1866.24	1804.032
STS-48	OC-48	STM-16	2488.32	2405.376
STS-96	OC-96	STM-32	4976.64	4810.176
STS-192	OC-192	STM-64	9953.28	9620.928

STS-3c Frame Format

- □ STS-3c is similar to STM-1
- \square 125 µs = 2430 bytes at 155.54 Mbps



Path Overhead

□ Note: All sizes are multiples of 3

Scrambling

- SONET uses NRZ coding.1 = Light On, 0 = Light Off.
- \square Too many 1's or 0's \Rightarrow Loss of bit clocking information
- □ All bytes (except some overhead bytes) are scrambled
- □ Polynomial $1 + x^6 + x^7$ with a seed of 11111111 is used to generate a pseudo-random sequence, which is XOR'ed to incoming bits.
 - 1111 1110-0000 0100-0001 ... 010
- ☐ If user data is identical to (or complement of) the pseudo-random sequence, the result will be all 0's or 1's.

Automatic Protection Switching

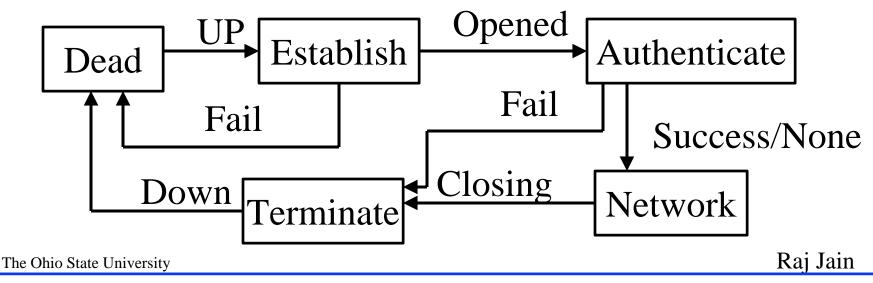
- 100 μs or more is "loss of signal"
 2.3 μs or less is not "loss of signal"
 In-between is up to implementations
- Most implementations use 13-27 μ s ⇒ Higher speed lines ⇒ maintain sync for more bits
- □ APS allows switching circuits on fault
- May take up to 50 ms to complete
- Wastes entire links as standby.
- Protection by routers works faster than by SONET

HDLC Family

- Synchronous Data Link Control (SDLC): IBM
- □ High-Level Data Link Control (HDLC): ISO
- □ Link Access Procedure-Balanced (LAPB): X.25
- Link Access Procedure for the D channel (LAPD): ISDN
- □ Link Access Procedure for modems (LAPM): V.42
- □ Link Access Procedure for half-duplex links (LAPX): Teletex
- □ Point-to-Point Protocol (PPP): Internet
- Logical Link Control (LLC): IEEE
- Advanced Data Communications Control Procedures (ADCCP): ANSI
- V.120 and Frame relay also use HDLC

PPP: Introduction

- Point-to-point Protocol
- Originally for User-network connection
- Now being used for router-router connection
- Three Components: Data encaptulation, Link Control Protocol (LCP), Network Control Protocols (NCP)



PPP Procedures

- □ Typical connection setup:
 - Home PC Modem calls Internet
 Provider's router: sets up physical link
 - PC sends series of LCP packets
 - Select PPP (data link) parameters
 - Authenticate
 - PC sends series of NCP packets
 - Select network parametersE.g., Get dynamic IP address
- Transfer IP packets

PPP in HDLC-Like Framing

Flag Address Co.

Control | F

Protocol

01111110 11111111 00000011

Info Padding CRC Flag

- □ Flag = 0111 1110 = 7E
- □ Byte Stuffing:

 $7E \Rightarrow 7D 5E$

 $7D \Rightarrow 7D 5D$

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Framing (Cont)

- \square Address=FF \Rightarrow All stations
- □ Control=03 ⇒ Unnumbered
 Poll/final = command/response = 0 ⇒ Response
- □ Protocol = 8/16 bits. lsb=1 of LSB \Rightarrow End of address All protocols are odd and lsb of MSB = 0
- Packets may be padded up to MRU.

 Maximum receive unit = 1500 default
- □ 16-bit FCS default 32-bit FCS can be negotiated using LCP
- HDLC Shared zero mode:

O111111101111110 = Flag-Flag. Not used in PPP_{Raj Jain}

LCP Config Options

- Maximum Receive Unit
- □ Authentication Protocol: C0 23 ⇒ Password
 C2 23 ⇒ Challenge Handshake
- \square Quality Protocol: C025 \Rightarrow Will expect link reports
- Magic Number: To related responses with requests Randomly number in sequence of the request Helps detect looped back links
- Protocol Field Compression:Only one byte is used even for 2-byte protocols
- □ Address and Control Field Compression:
 FF03 is not transmitted. CRC is on compressed frame.

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PPP over SONET

- □ In PPP/LCP [RFC 1619]:
 - Magic Number
 - No Address and Control field compression
 - No Protocol field compression
 - o 32-bit FCS
- ☐ In SONET:
 - No Scrambling of Payload
 - PATH Signal Label (C2) = 207 = CF
 - = Content type

SONET vs ATM

1. Overhead:

- SONET claimed to provide 25-30% higher throughput than ATM.
- IPOA encaptulation, AAL5 trailer, ATM cell headers eliminated in SONET
- 155.52 Mbps Link \Rightarrow 149.76 ATM \Rightarrow 135.63 ATM payload
- 9.5% more throughput (135.63 Mbps vs 149 Mbps)= 9 T1 Lines out of 96
- 6% for ABR flow control. Nothing for UBR/CBR/VBR.
- Signaling overhead for SVCs.

SONET vs ATM (Cont)

- 2. SONET Reliability through APS
 APS wastes entire links as standby.
 Long APS times can badly interact with routing
- 3. ATM provides multiservice integration
- 4. ATM provides traffic management (oversubscription)
- 5. SONET needs to be provisioned. ATM allows SVCs also.
- 6. ATM allows multiple secure VCs on the same physical interface.

SONET vs ATM (Cont)

- 7. SONET managed by TL-1 protocol. Will migrate to CMIP. IP and ATM can be managed by SNMP. Can't configure SONET equipment/ bandwidth from IP platform.
- 8. PPP byte stuffing create unpredictable traffic ⇒ QoS difficult
- 9. No Priorities or preemption in IP/PPP/SONET

 ⇒ QoS not feasible currently
- 10. PPP is a single-destination protocol. You can reach only one destination using one link. ATM is a multi-destination protocol.

SONET vs ATM (Cont)

SONET allows multiple destinations from one link using multiple OC-n frames but PPP cannot use this feature.

11. Multicast: No support in SONET.

Handled in IP.

Multicast over SONET being designed.

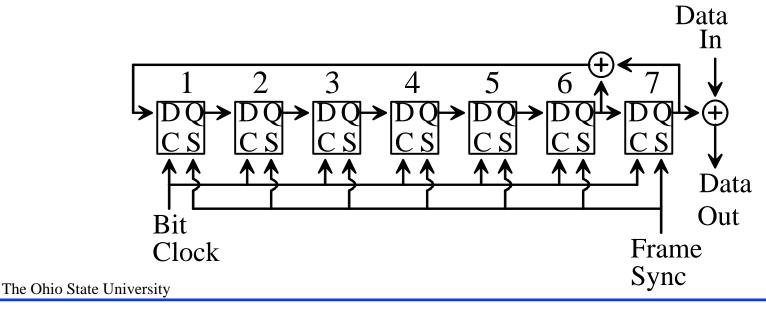
Multiple Access Protocol Over SONET (MAPOS)

- 12. Delay: Every hop of SONET introduces a 125-µs delay regardless of speed
 - ⇒ Cut through routing is difficult
- 13. SONET payload scrambling is an issue.

Scrambling: Introduction

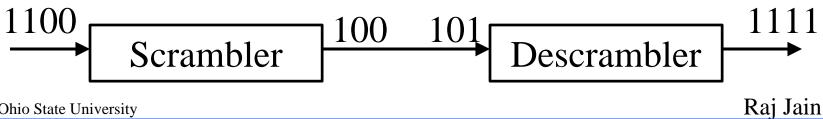
- 1. Add random sequence
- 2. Divide by a number and send quotient. Similar to CRC.

Both implemented by shift-registers. Analyzed using polynomials. $1+x^6+x^7$



Scrambling (Cont)

- □ Set-Reset Synchronous scrambler: Add a fixed random bit pattern. Need to tell where to start adding \Rightarrow Need to synchronize.
- Self-synchronous scrambler: Divide by a fixed number. No need for synchronization. Errors multiply.
- \square Example: Send 12 using divider 3 \Rightarrow Send 4. Received $5 \Rightarrow 15$.



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Payload Scrambling Issue

- 21 1500-byte datagrams will ensure
 2080 bits of 0's/1's (13 μs at STS-3c)
 resulting in Loss of signal, framing, and Sync
 [T1X1.5/97-134, 97-130]
- □ Standard requires 2.3-100 µs LOS.

 Most interfaces are on the low end.

 Most interfaces can't keep clock sync after 80 bits
- □ Carriers tariffs based on failures and errors guarantees
 ⇒ Customer can cause excessive failures and no way for carriers to trace it.
- □ A single packet can disrupt a large number of users.
- \square APS is triggered \Rightarrow Disruption could last up to 50 ms.

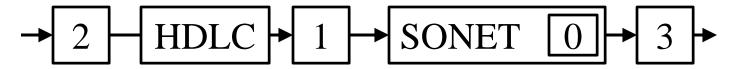
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Scrambling: Solutions

- 1. ANSI T1X1.5+IETF recommend using 1+x⁴³ for PPP over SONET for STS-1 through STS-48. Higher or lower rates require further study.
- A path signal label different from 207 will be used to differentiate scrambled and non-scrambled payloads.
- □ Self-synchronous scrambler \Rightarrow error-multiplying. 1-bit error on the line \Rightarrow 2-bit errors in packet
- □ Some error patterns detectable w/o scrambler are undetectable with scrambler
- □ FCS bit ordering (lsb) and scrambler bit ordering (msb) also have some effect.

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Solutions (Cont)



- 2. Scramble PPP before HDLC framing
 - ⇒ Requires disabling errored HDLC frame discard. Does not protect against framer errors.
- 3. Scramble the SONET scrambler output.
- 4. Use $1+x^2+x^{19}+x^{21}+x^{40}$ set-reset frame synchronous scrambler
- 5. Avoid long sequences of zeros in the SONET scrambler output by pattern matching HDLC packet and byte-stuffing.



- □ IP over SONET = IP over PPP in HDLC-like framing over SONET/SDH
- SONET does not provide QoS, Dynamic bandwidth (SVCs), QoS multiplexing, traffic management
- Payload scrambling is a hot issue

References

- □ For a detailed list of references, see http://www.cis.ohio-state.edu/~jain/refs/snt_refs.htm
- □ RFC 1619, PPP over SONET/SDH,
- □ RFC 1662, PPP in HDLC-like Framing
- □ RFC 1661, The Point-to-Point Protocol (PPP)
- □ "PPP Over SONET Mapping", 10/23/1997, draftallen-pppsonet-mapping-00.txt
- □ "PPP over SONET/SDH", 10/16/1997, draft-ietf-pppext-pppsonet-scrambler-00.txt
- □ "PPP over SONET/SDH", 11/17/1997, draft-ietf-pppext-sonet-ds-00.txt

IP over SONET: Products

CISCO:

- Cisco 7505, Cisco 7507, Cisco 7513: up to 6 OC-3
 Cisco 12004,8,12: 5-60 Gbps, 12-44 OC-3, 3-11
 OC-12
- 12012 has 12 chassis slots. Up to 5 can be configured as fabric (15-60 Gbps)
- Up to 11 can be configured for line cards
- → Line cards: 4-port OC-3 or 1-port OC-12
- → 4-port OC-12 and 1-port OC-48 planned.

Products (Cont)

- IP, IPX, DECnet over SONET
- Payload scrambling
- Automatic protection switching
- Monitoring and fault detection
- Line rate forwarding
- Carrier Class: redundant processors, fabric, power, fans, line cards, hot swap
- o Ref:

http://cio.cisco.com/warp/public/733/12000/gsrfs_ds.htm http://www.cisco.com/warp/public/733/12000/gspos_an.htm

Products (Cont)

- □ Ascend: GRF 400 IP switches
 - **2.8** Mpps
 - IP forwarding media card
 - 120 kpps
 - Hot swap, Automatic Protection switching, CISCO compatible HDLC
- Essential Communications: ERF IP Gateway (Cross-point Switch) [Same as Ascend?]
 - ERF-400 up to 4 Gbps, 2.8 Mpps
 - ERF-1600 up to 16 Gbps, 10 Mpps
 - OC-3c and OC-12c

Abbreviations

ABR Available Bit Rate

ACCM Async Control Character Map

ANSI American National Standards Institute

APS Automatic Protection Switching

ATM Asynchronous Transfer Mode

CCITT Consultative Committee

CMIP Common Management Information Protocol

CMISE Common Management Information Service Element

COBS Consistent overhead byte stuffing

CRC Cyclic Redundancy Check

FCS Frame Check Sequence

FDDI Fiber Distributed Data Interface

HDLC High-Level Data Link Control

IP Internet Protocol

IPCP IP Control Protocol

IPOA IP over ATM

IPV6CP IPv6 Control Protocol

ISDN Integrated Service Digital Network

L2TP Layer Two Tunneling Protocol

LCP Link Control Protocol

LoS Loss of Signal

lsb Least significant bit

LSB Least Significant Byte

MPOA Multiprotocol over ATM

MPV Maximum pad value

MRU Maximum Receive Unit

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msb Most significant bit

MSB Most Significant Byte

MSP Multiplex Section Protection Function

MTU Maximum transmission unit

NCP Network Control Protocols

NRZ Non-return to Zero

OC Optical Carrier

PPP Point-to-Point Protocol

PPTP Point-to-point Tunneling Protocol

QoS Quality of Service

RFC Request for Comments

SNMP Simple Network Management Protocol

SONET Synchronous Optical Network

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SPE Synchronous Payload Envelop

SVC Switched Virtual Circuit

TCP Transmission Control Protocol

XOR Exclusive-OR

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